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GoFFish: A Sub-Graph Centric Framework for Large-Scale Graph Analytics

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GoFFish Overview

- Scalable platform for *large scale graph analytics* in <u>commodity clusters and cloud</u>.
- Simple yet scalable programming abstraction for large scale graph processing
- More than 10x improvements in some graph algorithms over traditional graph processing systems
- Components
 - GoFS : Sub-graph centric *Distributed Graph Storage*
 - Gopher : Distributed BSP-based programming abstraction for Large Scale graph analytics framework

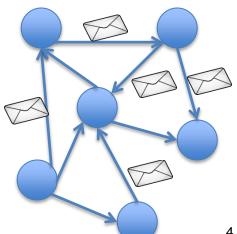


Vertex centric programming model

- Iterative vertex centric programming model based on Bulk synchronous parallel model.
- In each iteration vertex can
 - 1. Receive messages sent to it in previous iteration
 - 2. Send messages to other vertices
 - 3. Modify its own state
- Vertex centric programming mode- *Think like a vertex*
- Very simple,

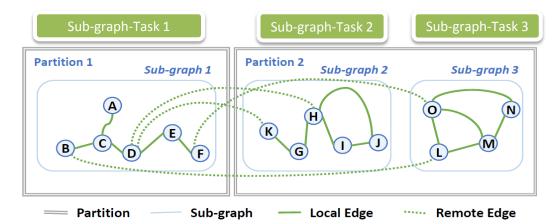
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- High communication overhead
- Pregel, <u>http://giraph.apache.org/</u>



Subgraph Centric Programming Model

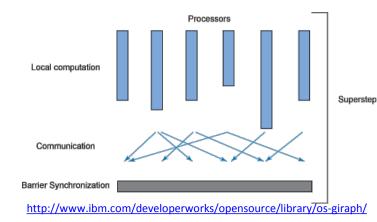
- Partition graph in to set of sets of vertices and edges
- *Sub-graphs* Connected components in a partition.
- User logic operates on a sub graph
 - Independent unit of computation
- Resource allocation
 - − Single Partition \rightarrow Single Machine | Single Sub-graph \rightarrow Single CPU
- Data loading
 - Entire partition is loaded into memory before computation
 - Tasks retain sub-graphs in memory within the task scope

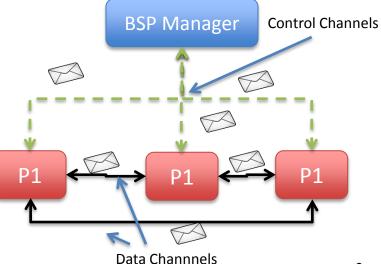




Sub-graph Centric Programming Model

- <u>Algorithm</u>: Seqⁿ of super-steps separated by global barrier synchronization
- Super Step i:
 - 1. Sub-graphs compute in parallel
 - 2. Receive messages sent to it in super step i-1
 - 3. Execute same user defined function
 - 4. Send messages to other sub graphs (to be received in super step i+1)
 - 5. Can vote to halt : I'm done / de-activate
- Global Vote to Halt check
 - termination \rightarrow all sub-graphs voted to halt
- Active sub graphs participate in every computation
- De-activated sub-graphs will not get executed/activated unless it get new messages





Sub-graph Centric Programming Model Example

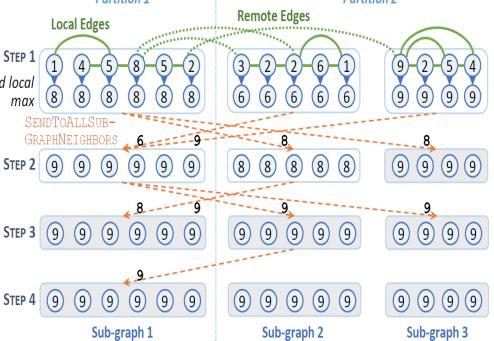
- Algorithm Max vertex value
- Input Connected graph with different vertex values {n₁,n₂.... n_i}
- Output Each vertex with the Max ({n₁,n₂..., n_i})

```
Algorithm 1 Max Vertex using Vertex Centric Model
 1: procedure COMPUTE(Vertex myVertex, Iterator(Message) M)
       hasChanged = (superstep == 1) ? true : false
 2:
       while M.hasNext do
                                   ▶ Update to max message value
 3:
           Message m \leftarrow M.next
 4:
          if m.value > myVertex.value then
 5:
                                                                                Partition 1
                                                                                                                     Partition 2
              myVertex.value \leftarrow m.value
 6:
                                                                                                     Remote Edges
                                                                          Local Edges
              hasChanged = true
 7:
                                                                   STEP 1
       if hasChanged then
                                  ▶ Send message to neighbors
 8:
                                                                                                                            9
           SENDTOALLNEIGHBORS(myVertex.value)
                                                                Find local
 9:
                                                                              8
       else
10:
                                                                    тах
           VOTETOHALT()
11:
```

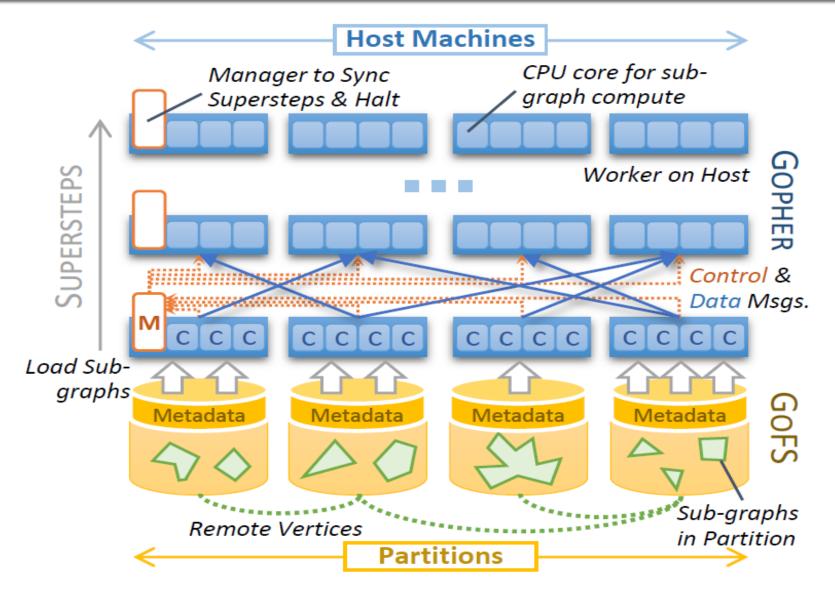
- Compared to Vertex centric model
 - Less communication
 - Fast convergence

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Goffish Architecture





Performance Comparisons – Experimental Setup

- Cluster
 - 12 Nodes, 8-Core Intel Xeon CPU (each)
 - 16GB RAM (each), 1TB HDD (each)
- Network

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- Gigabit Ethernet
- Apache Giraph latest version from trunk
 - Includes Performance improvements from Facebook !

Data Set	Vertices	Edges	Diameter
RN: California Road Network	1,965,206	2,766,607	849
TR: Internet path graph from Tracesroutes	19,442,778	22,782,842	25
LJ: LiveJournal Social Network	4,847,571	68,475,391	10
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Runtime Performance Comparison with Vertex Centric Model

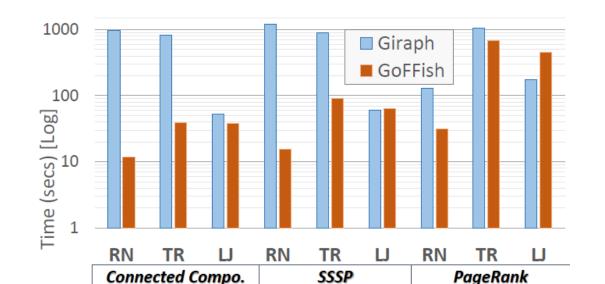
- Connected components
 - 81x improvement using California Road Network (RN) dataset
 - 21x improvement using a Trace route path network of a CDN. (TR) dataset
- Single Source Shortest Path

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- 32x improvement using RN dataset
- 10x improvement using TR dataset

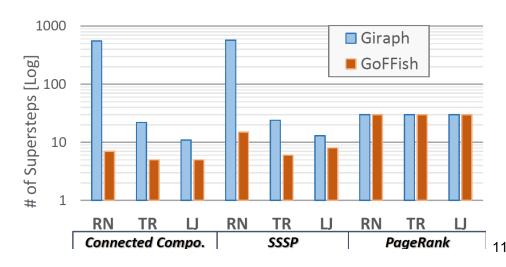
Page Rank

4x improvement using RN dataset 1.5x improvement using TR dataset Not an ideal algorithm for sub-graph centric programming model



Super Step Reduction Comparison with Vertex Centric Model

- Connected components
 - ~79x reduction using RN dataset.
 - ~4x reduction using TR dataset
 - ~2x reduction using Live Journal dataset (LJ)
- Single source shortest path
 - ~38x reduction using RN data set
 - 4x reduction using TR dataset
 - ~1.6x reduction using Live journal dataset





Conclusion

- Introduced a *sub-graph centric programming abstraction* for large scale graph analytics on distributed systems
 - Simple
 - Enable using shared memory algorithms at sub-graph level.
- Sub-graph centric algorithms and performance results
 - Connected Components
 - Single Source Shortest Path
 - Page Rank
- Issues and Future work
 - Sub-graph aware partitioning
 - Sub-graph centric algorithms
- Try now
 - <u>https://github.com/usc-cloud/goffish</u>



Thank you!



http://thesciencepresenter.wordpress.com/category/behaviour-management/

